

Bachelor's Thesis

Efficient Synchronization of Linux Memory Regions over a Network

A Comparative Study and Implementation

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r3map



Loophole Labs

Introduction

Today's Technological Landscape

What if there were a better way?

Thesis Structure

1. Introduction to Base Technologies
2. Overview of Access Methods
3. Implementing Select Access Methods
4. Analyzing Performance Benchmarks
5. Discussion of Benchmarks and Technologies
6. Conclusion and Future Outlook

Presentation Structure

1. Introduction
2. Methods
3. Optimizations
4. Results and Discussion
5. Implemented Use Cases
6. Future Use Cases
7. Conclusion

Access Methods

Userfaults

Principle of Locality

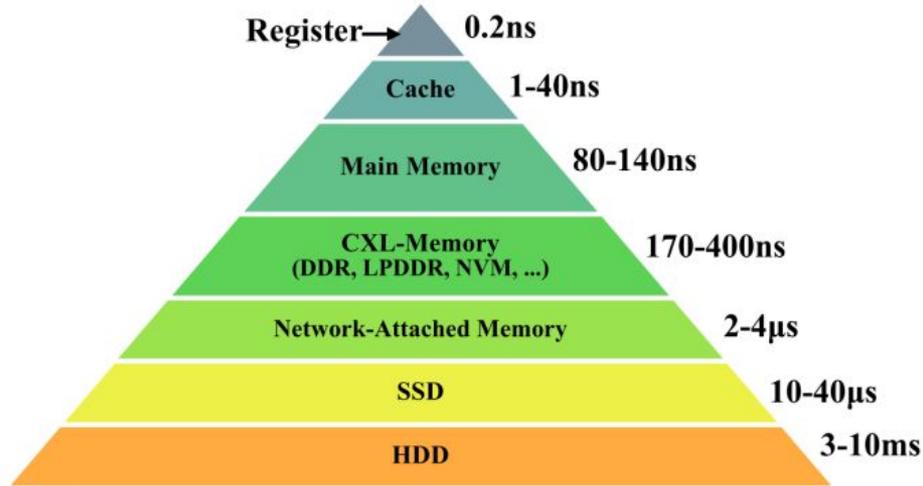


Figure 1: Latencies for different memory technologies showing, from lowest to highest latency, registers, cache, main memory, CXL memory, network-attached memory, SSDs and HDDs [12]

Memory Hierarchy

Page Faults

An Overview of Userfaults

```
1 // Creating the `userfaultfd` API
2 uffd, _, errno := syscall.Syscall(constants.NR_userfaultfd, 0, 0, 0)
3
4 uffdioAPI := constants.NewUffdioAPI(
5     constants.UFFD_API,
6     0,
7 )
8 // ...
9
10 // Registering a region
11 uffdioRegister := constants.NewUffdioRegister(
12     constants.CULong(start),
13     constants.CULong(1),
14     constants.UFFDIO_REGISTER_MODE_MISSING,
15 )
16 // ...
17 syscall.Syscall(
18     syscall.SYS_IOCTL,
19     uffd,
20     constants.UFFDIO_REGISTER,
21     uintptr(unsafe.Pointer(&uffdioRegister))
22 )
```

Implementing Userfaults

```
1 func (a abcReader) ReadAt(p []byte, off int64) (n int, err error) {
2     n = copy(p, bytes.Repeat([]byte{'A' + byte(off%20)}, len(p)))
3
4     return n, nil
5 }
```

```
1 f, err := os.OpenFile(*file, os.O_RDONLY, os.ModePerm)
2 b, uffd, start, err := mapper.Register(int(s.Size()))
3 mapper.Handle(uffd, start, f)
```

```
1 // ...
2 f, err := mc.GetObject(ctx, *s3BucketName, *s3ObjectName, minio.GetObjectOptions{})
3 b, uffd, start, err := mapper.Register(int(s.Size()))
4 mapper.Handle(uffd, start, f)
```

Implementing Userfaults

File-Based Synchronization

mmap

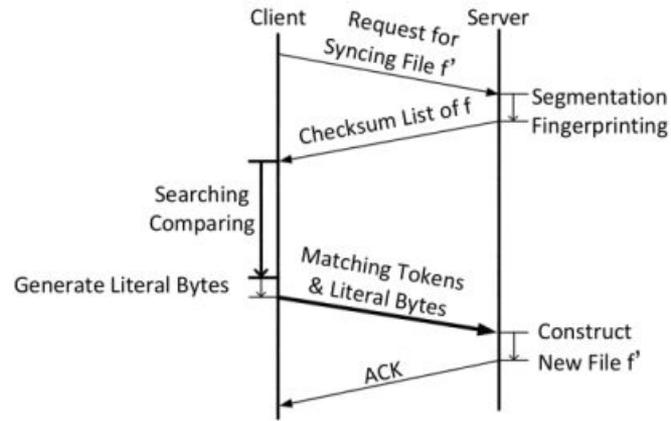


Figure 2: Design flow chart of WebRsync, showing the messages sent between and operations done for server and client in a single synchronization cycle[25]

Delta Synchronization

How can we use mmap and Delta Synchronization?

Caching Restrictions

Detecting File Changes

Speeding Up Hashing

Protocol & Multiplexer Hub

```

1 // Receiving remote hashes
2 remoteHashes := []string{}
3 utils.DecodeJSONFixedLength(conn, &remoteHashes)
4 // ...
5
6 // Calculating the hashes
7 localHashes, cutoff, err := GetHashesForBlocks(parallel, path, blocksize)
8
9 // Comparing the hashes
10 blocksToSend := []int64{}
11 for i, localHash := range localHashes {
12     // ...
13     if localHash != remoteHashes[i] {
14         blocksToSend = append(blocksToSend, j)
15     }
16     continue
17 }
18 }
19
20 // Sending the non-matching hashes
21 utils.EncodeJSONFixedLength(conn, blocksToSend)

```

```

1 // The lock and semaphore
2 var wg sync.WaitGroup
3 wg.Add(int(blocks))
4
5 lock := semaphore.NewWeighted(parallel)
6
7 // ...
8
9 // Concurrent hash calculation
10 for i := int64(0); i < blocks; i++ {
11     j := i
12
13     go calculateHash(j)
14 }
15 wg.Wait()

```

```

1 // Local hash calculation
2 localHashes, _, err := GetHashesForBlocks(parallel, path, blocksize)
3 // Sending the hashes to the remote
4 // Receiving the remote hashes and the truncation request
5 blocksToFetch := []int64{}
6 utils.DecodeJSONFixedLength(conn, &blocksToFetch)
7 // ...
8 cutoff := int64(0)
9 utils.DecodeJSONFixedLength(conn, &cutoff)

```

```

1 case "src-control":
2     // Decoding the file name
3     file := ""
4     utils.DecodeJSONFixedLength(conn, &file)
5     // ...
6
7     syncerSrcControlConns[file] = conn
8
9     syncerSrcControlConnsBroadcaster.Broadcast(file)
10    // ...
11 case "dst-control":
12     var wg sync.WaitGroup
13     wg.Add(1)
14
15     go func() {
16         // Subscription to send all future file names
17         l := syncerSrcControlConnsBroadcaster.Listener(0)
18
19         for file := range l.Ch() {
20             utils.EncodeJSONFixedLength(conn, file)
21             // ...
22         }
23     }()
24
25     // Sending the previously known file names
26     for file := range syncerSrcControlConns {
27         utils.EncodeJSONFixedLength(conn, file)
28         // ...
29     }
30
31     wg.Wait()

```

Delta Synchronization

Limitations

FUSE

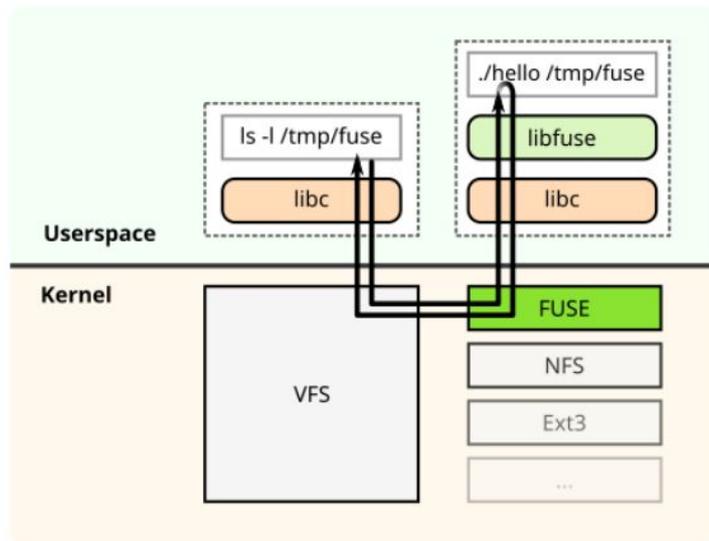


Figure 3: Structural diagram of FUSE, showing the user space components handled by the C library and the FUSE library as well as the kernel components such as the Linux VFS and the FUSE kernel module[27]

What is FUSE?

```
1 static int example_getattr(const char *path, struct stat *stbuf,  
2 struct fuse_file_info *fi);
```

```
1 static int example_readdir(const char *path, void *buf, fuse_fill_dir_t filler,  
2 off_t offset, struct fuse_file_info *fi,  
3 enum fuse_readdir_flags flags);
```

```
1 static int example_open(const char *path, struct fuse_file_info *fi);
```

```
1 static int example_read(const char *path, char *buf, size_t size, off_t offset, struct  
fuse_file_info *fi);
```

FUSE Syscalls

FUSE for Memory Synchronization

Limitations

NBD

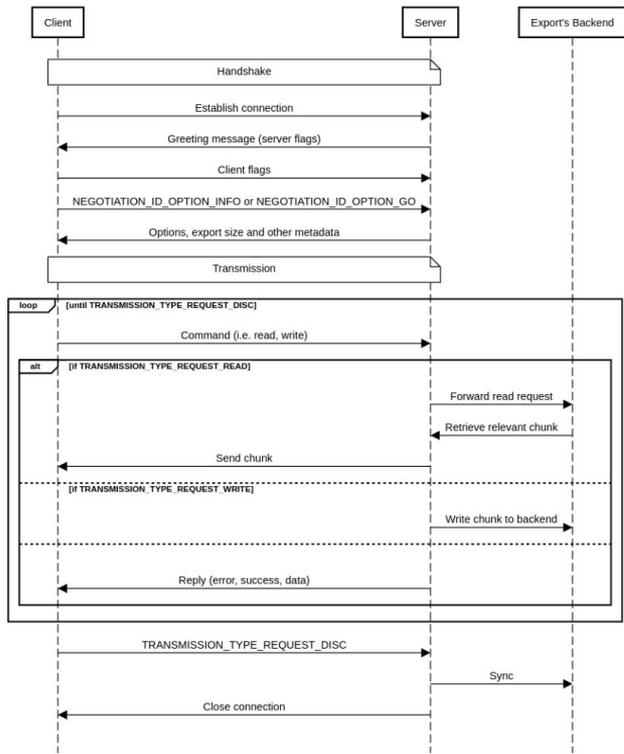


Figure 4: Sequence diagram of the baseline NBD protocol (simplified), showing the handshake, transmission and disconnect phases

NBD Protocol

NBD for Memory Synchronization



go-nbd

go-nbd

```
1 type Backend interface {
2     ReadAt(p []byte, off int64) (n int, err error)
3     WriteAt(p []byte, off int64) (n int, err error)
4     Size() (int64, error)
5     Sync() error
6 }
```

```
1 func Handle(conn net.Conn, exports []Export, options *Options) error
```

Server

```
1 func Connect(conn net.Conn, device *os.File, options *Options) error
```

```
1 // Connecting to `udev`  
2 udevConn.Connect(netlink.UdevEvent)  
3  
4 // Subscribing to events for the device name  
5 udevConn.Monitor(udevReadyCh, udevErrCh, &netlink.RuleDefinitions{  
6     Rules: []netlink.RuleDefinition{  
7         {  
8             Env: map[string]string{  
9                 "DEVNAME": device.Name(),  
10            },  
11        },  
12    },  
13 })  
14  
15 // Waiting for the device to become available  
16 go func() {  
17     // ...  
18     <-udevReadyCh  
19  
20     options.OnConnected()  
21 }()
```

Client

Combining Server and Client into Mounts

Managed Mounts

RTT, LAN and WAN



r3map

r3map

Why not just NBD?

Chunking

Background Pull and Push

```
1 type ReadWriterAt interface {  
2     ReadAt(p []byte, off int64) (n int, err error)  
3     WriteAt(p []byte, off int64) (n int, err error)  
4 }
```

Pipeline Stages

Pipeline Components

- ArbitraryReadWriterAt
- ChunkedReaderWriterAt
- SyncedReaderWriterAt: Background Pull and Push

Concurrent Initialization

```
1 type ManagedMountHooks struct {
2     OnBeforeSync func() error
3     OnBeforeClose func() error
4     OnChunkIsLocal func(off int64) error
5 }
```

Device Lifecycle

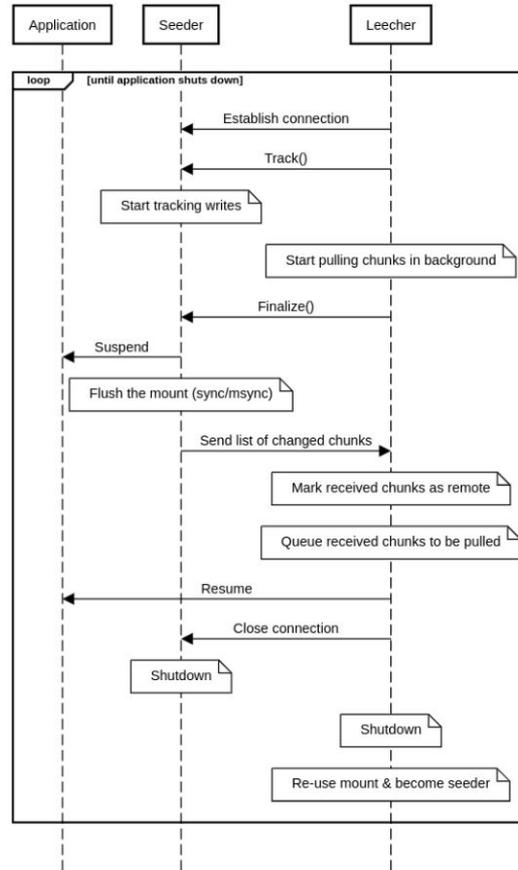
Live Migration

Pre-Copy Migration

Post-Copy Migration

Workload Analysis

Why a new API?



Migration Protocol

Figure 6: Sequence diagram of the migration protocol (simplified), showing the two protocol phases between the application that is being migrated, the seeder and the leecher components

```
1 type SeederRemote struct {
2     ReadAt func(context context.Context, length int, off int64) (r ReadAtResponse, err
      error)
3     Size    func(context context.Context) (int64, error)
4     Track   func(context context.Context) error
5     Sync    func(context context.Context) ([]int64, error)
6     Close   func(context context.Context) error
7 }
```

Seeder

```
1 // Suspends the remote application, flushes the mount and returns offsets that have been
   written too since `Track()`
2 dirtyOffsets, err := l.remote.Sync(l.ctx)
3
4 // Marks the chunks as remote, causing subsequent reads to pull them again
5 l.syncedReader.MarkAsRemote(dirtyOffsets)
6
7 // Schedules the chunks to be pulled in the background immediately
8 l.puller.Finalize(dirtyOffsets)
9
10 // Unlocks the local resource for reading
11 l.lockableReadWriteAt.Unlock()
```

Leecher

Optimizations

Pluggable Encryption, Authentication and Transport

Concurrent Backends

Remote Stores as Backends

```
1 func (b *RedisBackend) ReadAt(p []byte, off int64) (n int, err error) {
2     // Retrieve a key corresponding to the chunk from Redis
3     val, err := b.client.Get(b.ctx, strconv.FormatInt(off, 10)).Bytes()
4     // If a key does not exist, treat it as an empty chunk
5     if err == redis.Nil {
6         return len(p), nil
7     }
8     // ...
9 }
10
11 func (b *RedisBackend) WriteAt(p []byte, off int64) (n int, err error) {
12     // Store an offset as a key-value pair in Redis
13     b.client.Set(b.ctx, strconv.FormatInt(off, 10), p, 0)
14     // ...
15 }
```

Redis

```
1 func (b *S3Backend) ReadAt(p []byte, off int64) (n int, err error) {
2     // Receiving a chunk using Minio's S3 client
3     obj, err := b.client.GetObject(b.bucket, b.prefix+"-"+strconv.FormatInt(off, 10),
4         minio.GetObjectOptions{})
5     if err != nil {
6         // If an object is not found, it is treated as an empty chunk
7         if err.Error() == errNoSuchKey.Error() {
8             return len(p), nil
9         }
10        // ...
11    }
12    // ...
13 }
```

```
1 func (b *CassandraBackend) ReadAt(p []byte, off int64) (n int, err error) {
2     // Executing a select query for a specific chunk, then scanning it into a byte slice
3     var val []byte
4     if err := b.session.Query(`select data from `+b.table+` where key = ? limit 1`, b.
5         prefix+"-"+strconv.FormatInt(off, 10)).Scan(&val); err != nil {
6         if err == gocql.ErrNotFound {
7             return len(p), nil
8         }
9         return 0, err
10    }
11    // ...
12 }
13
14 func (b *CassandraBackend) WriteAt(p []byte, off int64) (n int, err error) {
15     // Upserting a row with a chunk's new content
16     b.session.Query(`insert into `+b.table+` (key, data) values (?, ?)`, b.prefix+"-"+
17         strconv.FormatInt(off, 10), p).Exec()
18     // ...
19 }
```

ScyllaDB

Dudirekta

gRPC

fRPC

Results and Discussion

Property	Value
Device Model	Dell XPS 9320
OS	Fedora release 38 (Thirty Eight) x86_64
Kernel	6.3.11-200.fc38.x86_64
CPU	12th Gen Intel i7-1280P (20) @ 4.700GHz
Memory	31687MiB LPDDR5, 6400 MT/s

Testing Environment

Access Methods

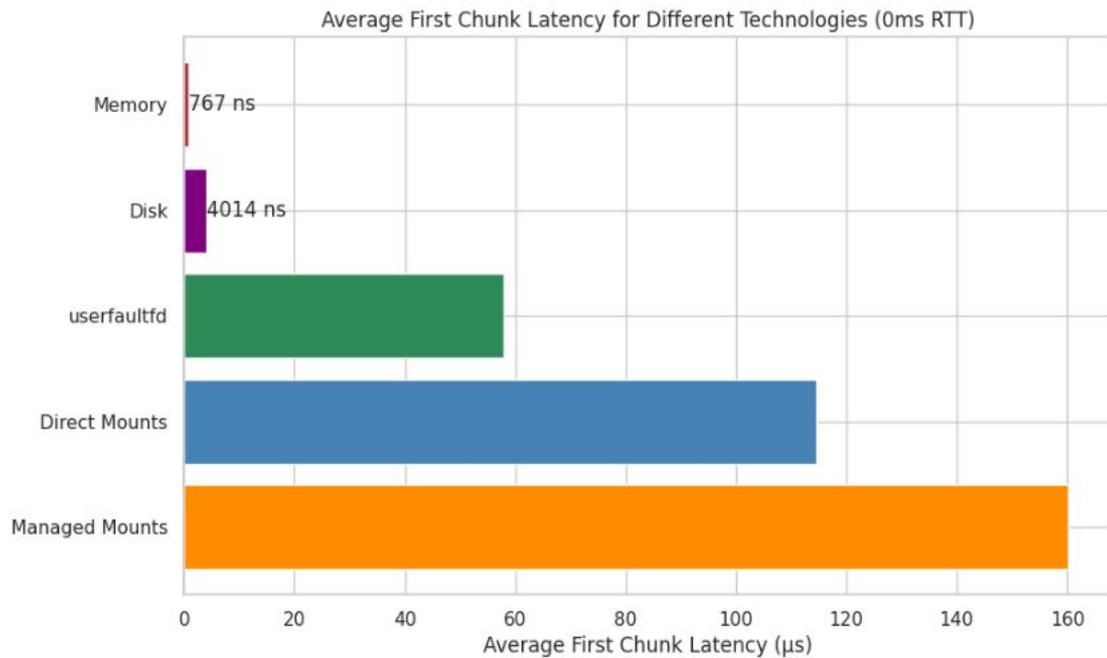


Figure 8: Average first chunk latency for different direct memory access, disk, userfaultfd, direct mounts and managed mounts (0ms RTT)

Latency

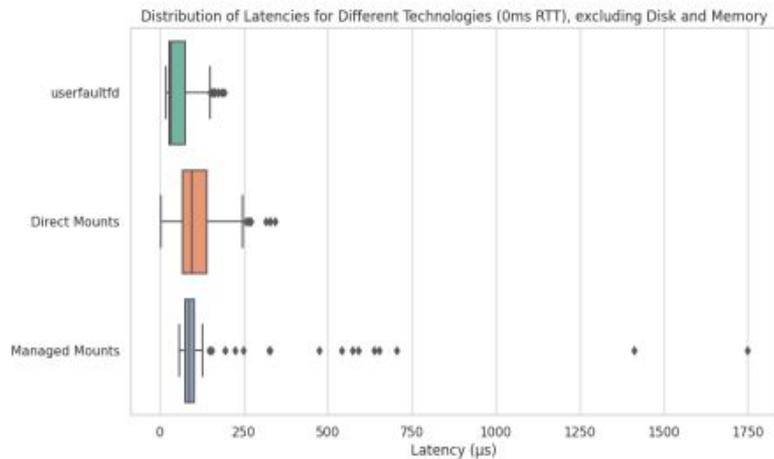


Figure 9: Box plot for the distribution of first chunk latency for userfaultfd, direct mounts and managed mounts (0ms RTT)

Latency

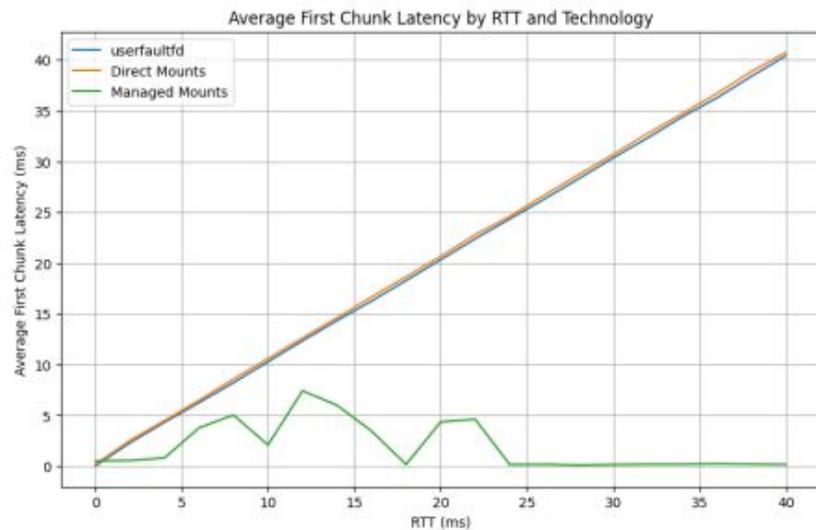


Figure 10: Average first chunk latency for userfaultfd, direct mounts and managed mounts by RTT

Latency

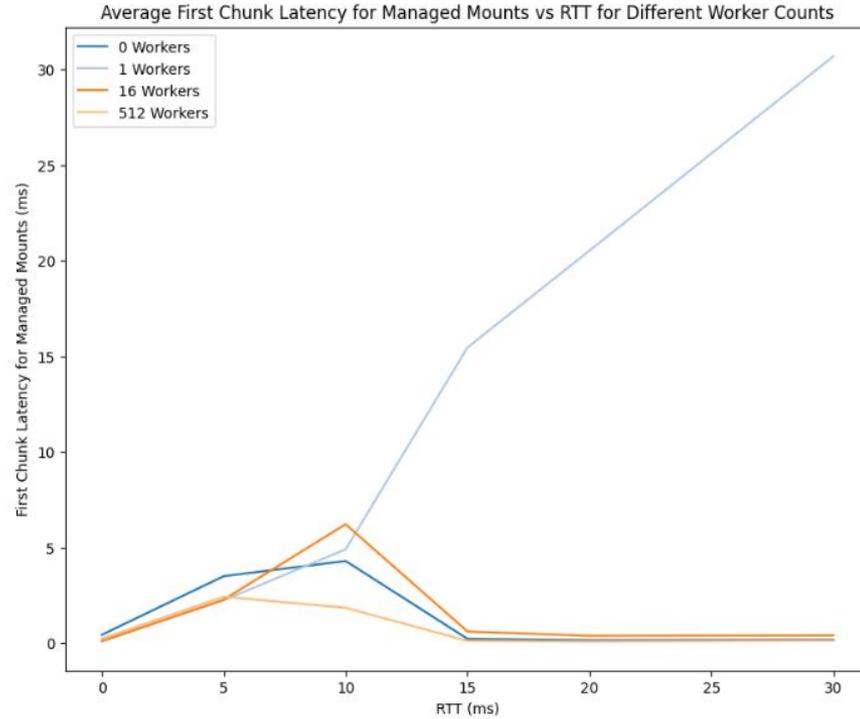


Figure 11: Average first chunk latency for managed workers with 0-512 workers by RTT

Latency

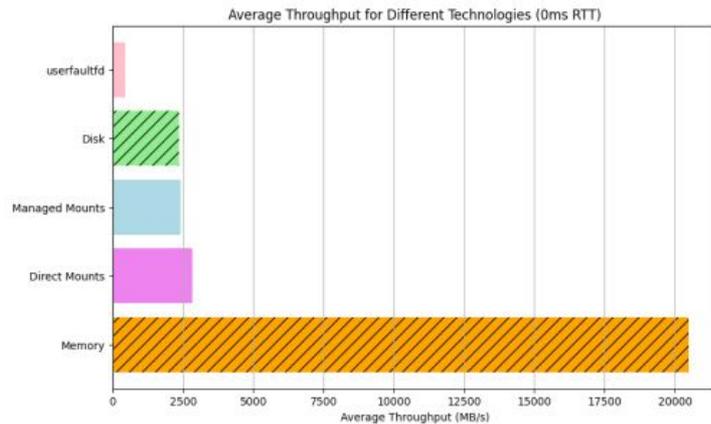


Figure 12: Average throughput for memory, disk, userfaultfd, direct mounts and managed mounts (0ms RTT)

Read Throughput

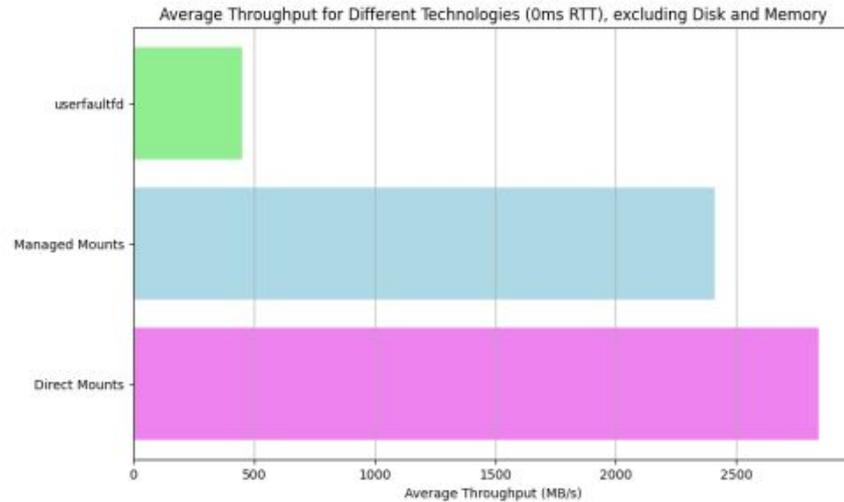


Figure 13: Average throughput for userfaultfd, direct mounts and managed mounts (0ms RTT)

Read Throughput

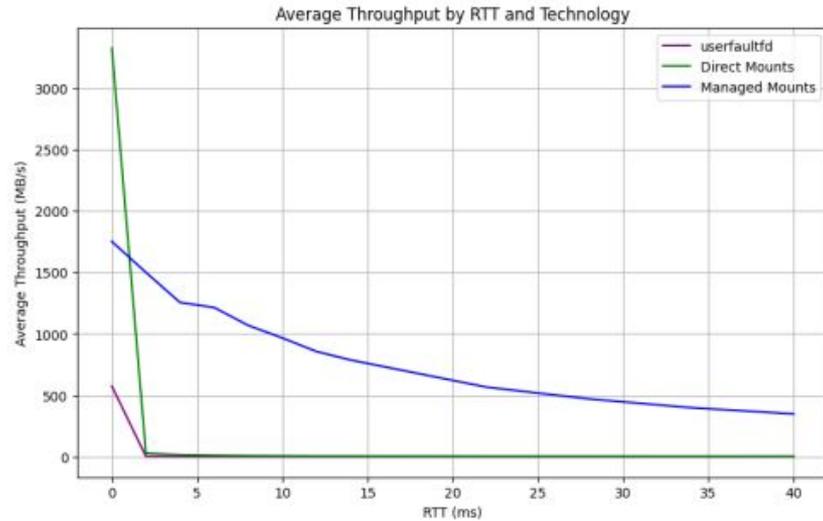


Figure 15: Average throughput for userfaultfd, direct mounts and managed mounts by RTT

Read Throughput

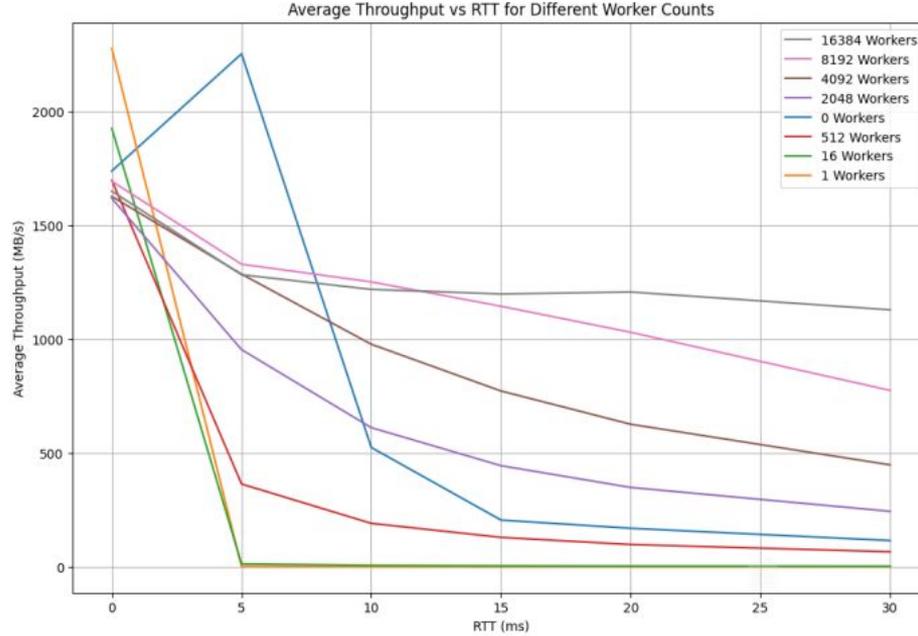


Figure 16: Average throughput for managed mounts with 0-16384 workers by RTT

Read Throughput

5.2.3 Write Throughput

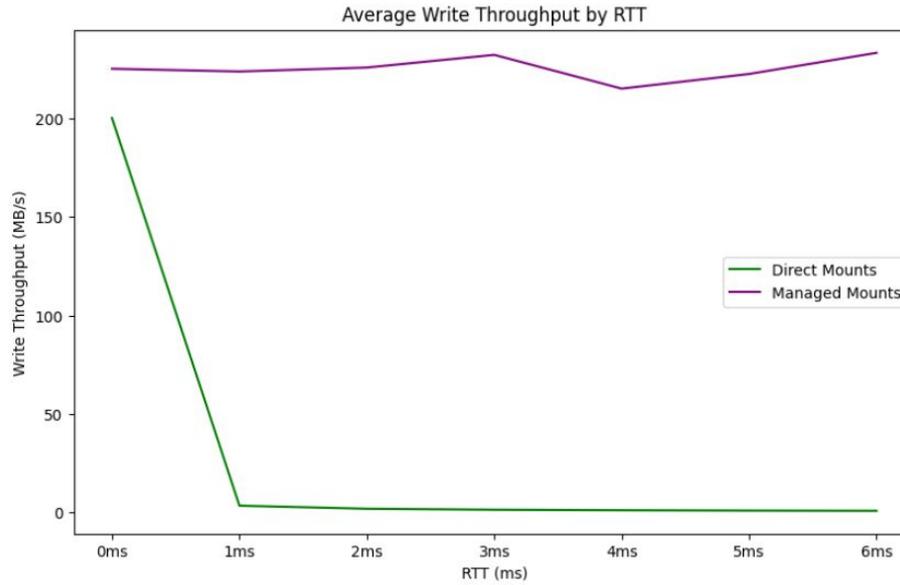


Figure 17: Average write throughput for direct and managed mounts by RTT

Write Throughput

Discussing Access Methods

Initialization

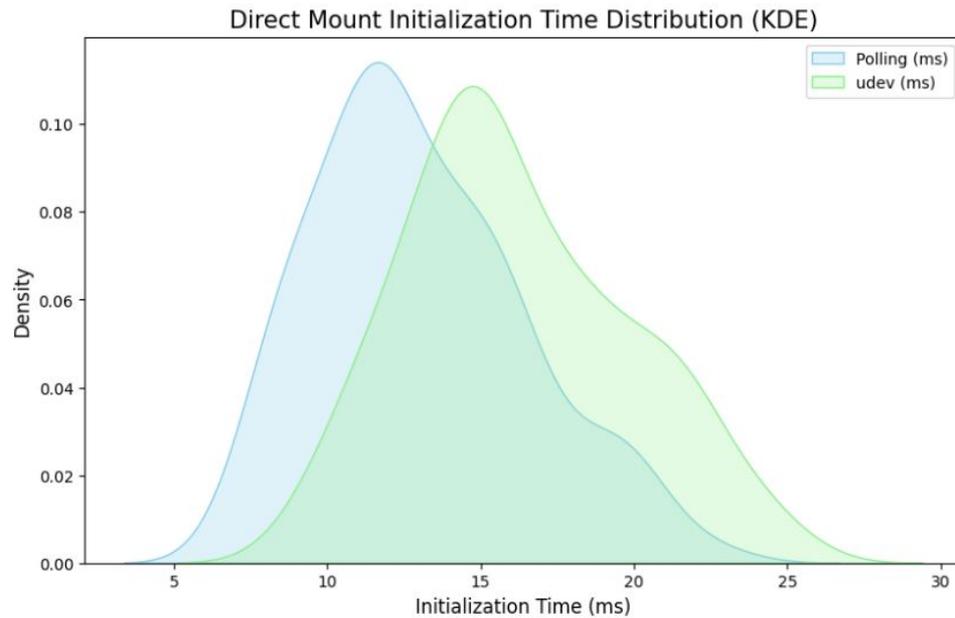


Figure 18: Kernel density estimation for the distribution of direct mount initialization time with polling vs. udev

Initialization

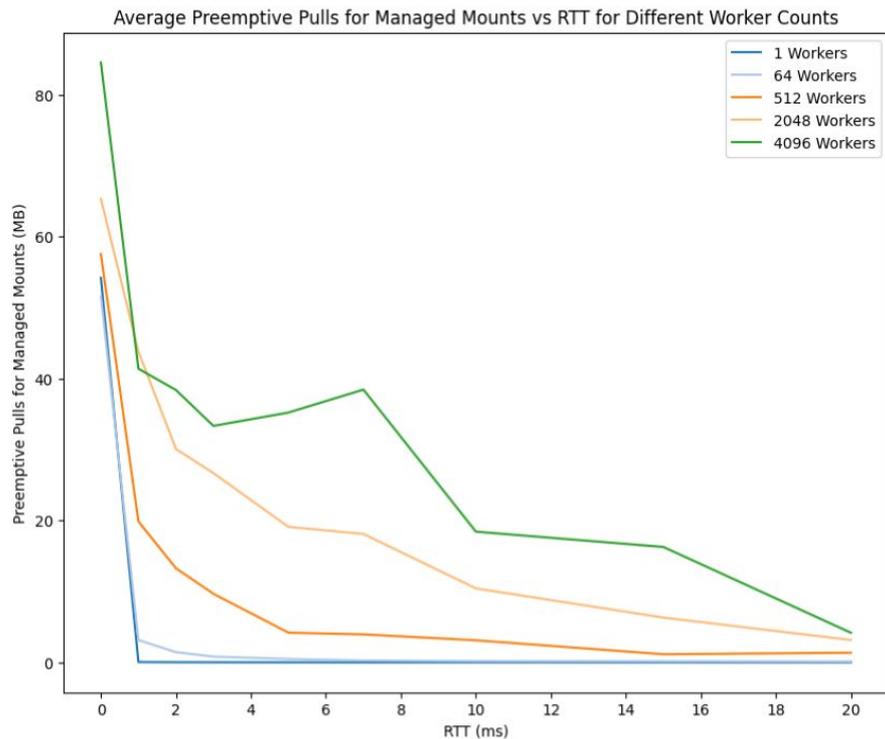


Figure 19: Amount of pre-emptively pulled data for managed mounts with 0-4096 workers by RTT

Initialization

Chunking

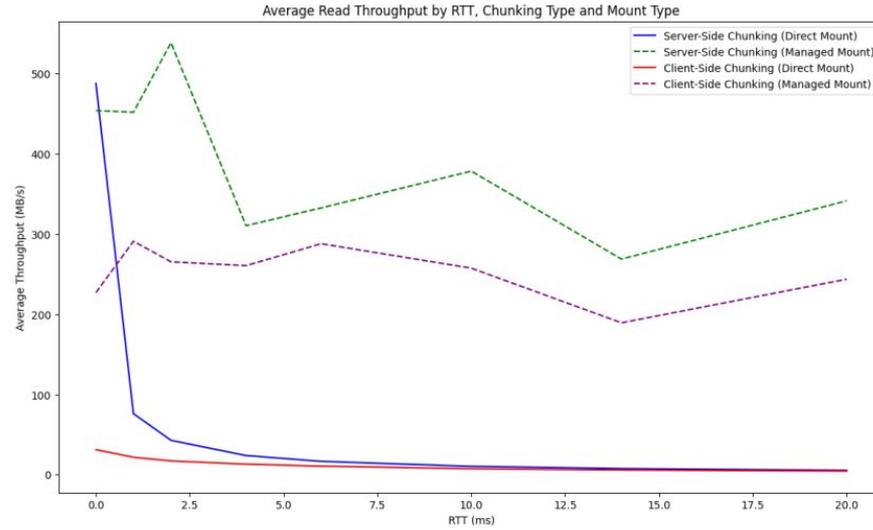


Figure 20: Average read throughput for server-side and client-side chunking, direct mounts and managed mounts by RTT

Read Throughput

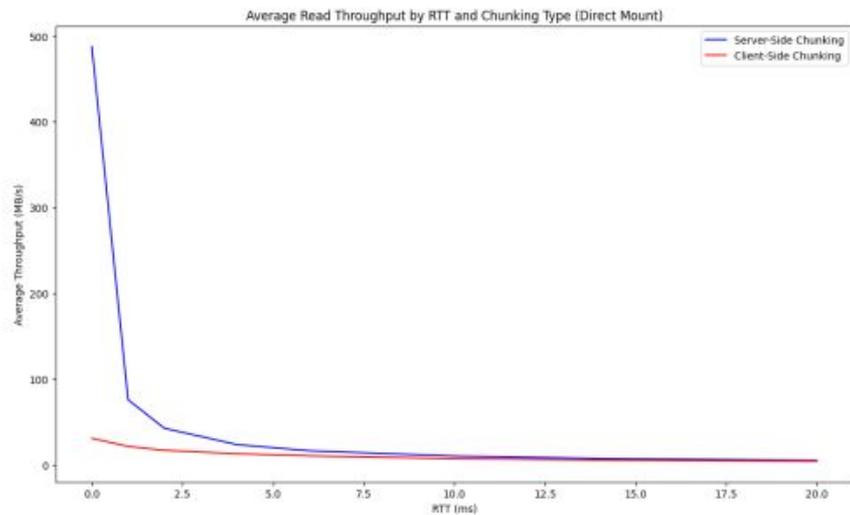


Figure 21: Average read throughput for server-side and client-side chunking with direct mounts by RTT

Read Throughput

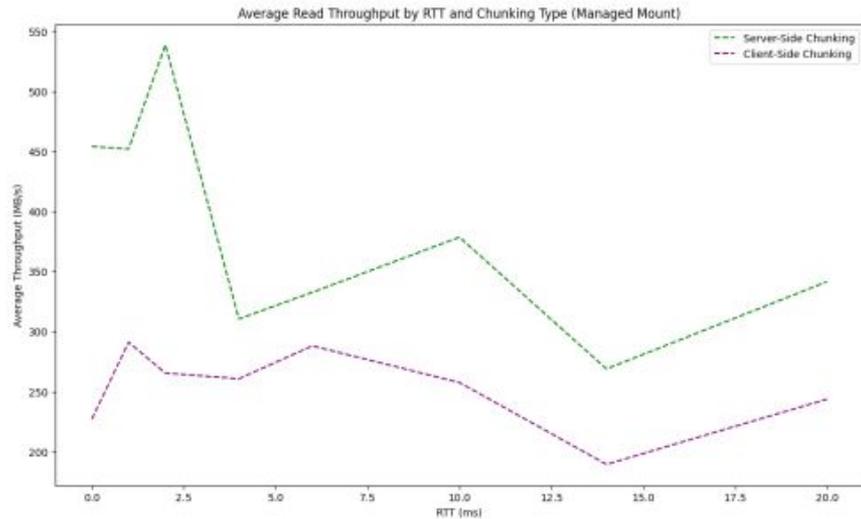


Figure 22: Average read throughput for server-side and client-side chunking with managed mounts by RTT

Read Throughput

RPC Frameworks

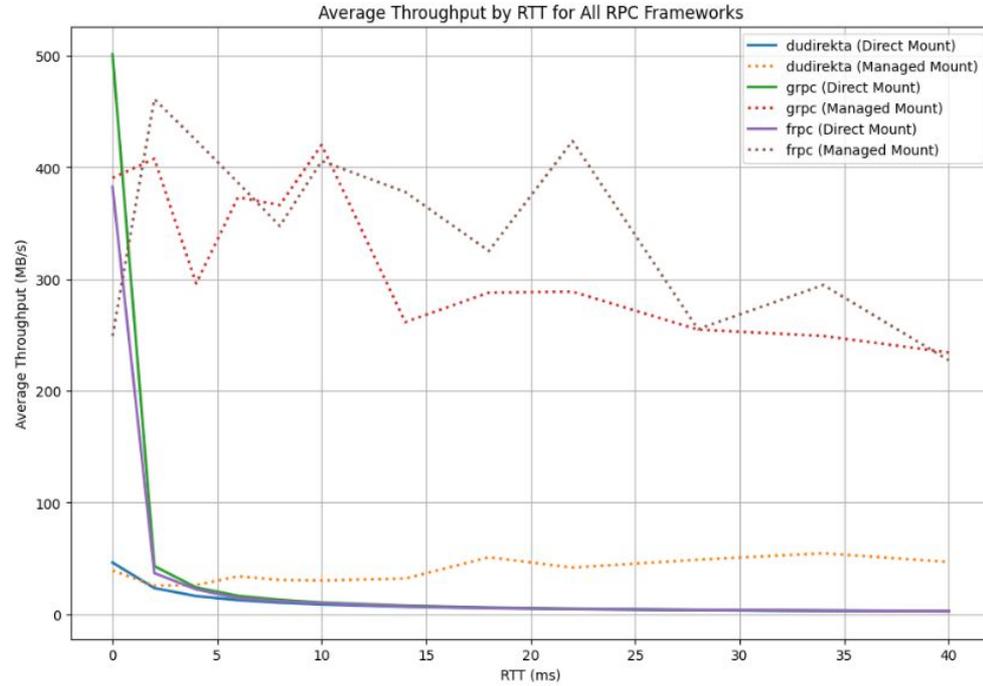


Figure 23: Average throughput by RTT for Dudirekta, gRPC and fRPC frameworks for direct and managed mounts

Throughput

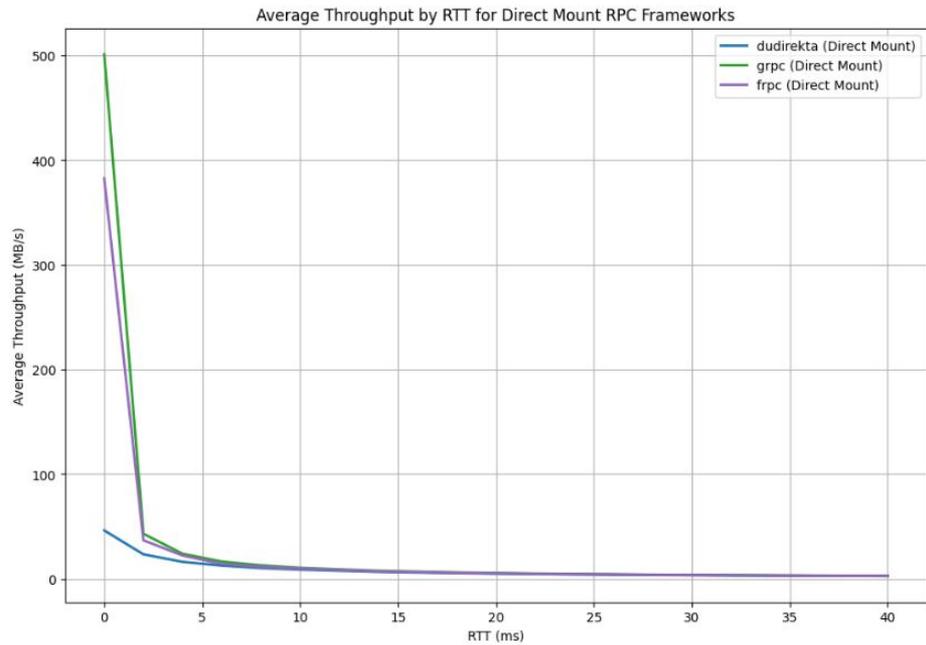


Figure 24: Average throughput by RTT for Dudirekta, gRPC and fRPC frameworks for direct mounts

Throughput

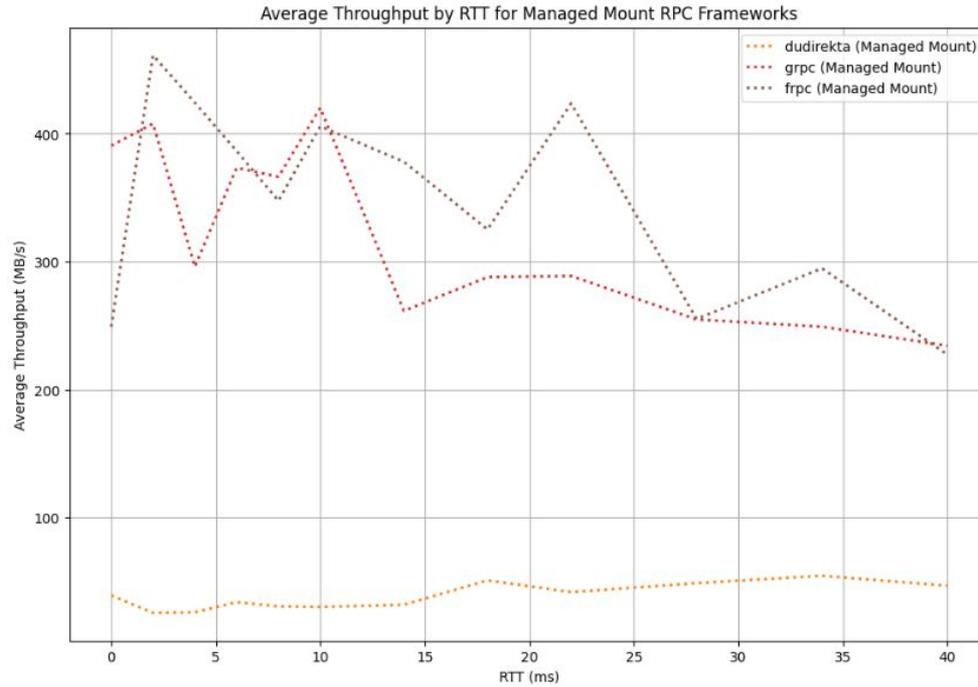


Figure 25: Average throughput by RTT for Dudirekta, gRPC and fRPC frameworks for managed mounts

Throughput

Discussing RPC Frameworks

Backends

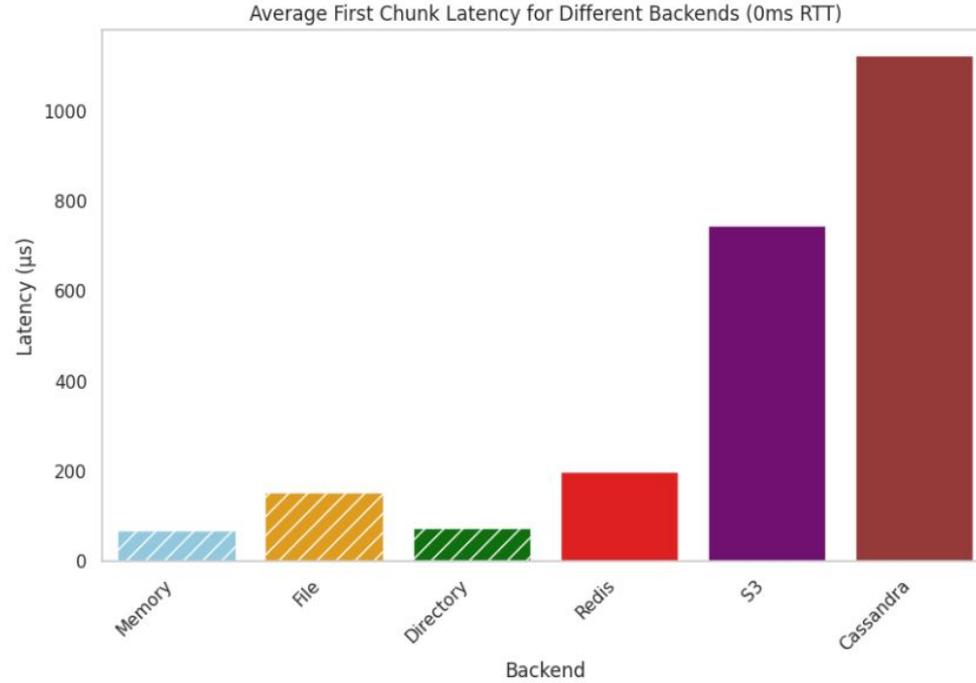


Figure 26: Average first chunk latency for memory, file, directory, Redis, S3 and ScyllaDB backends (0ms RTT)

Latency

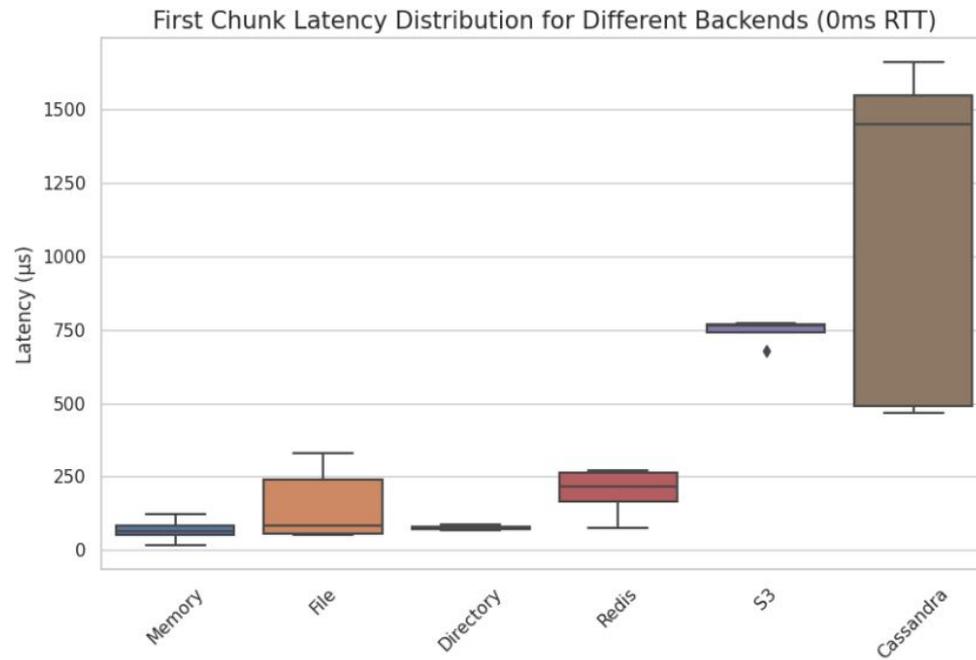


Figure 27: Box plot of first chunk latency distribution for memory, file, directory, Redis, S3 and ScyllaDB (0ms RTT)

Latency

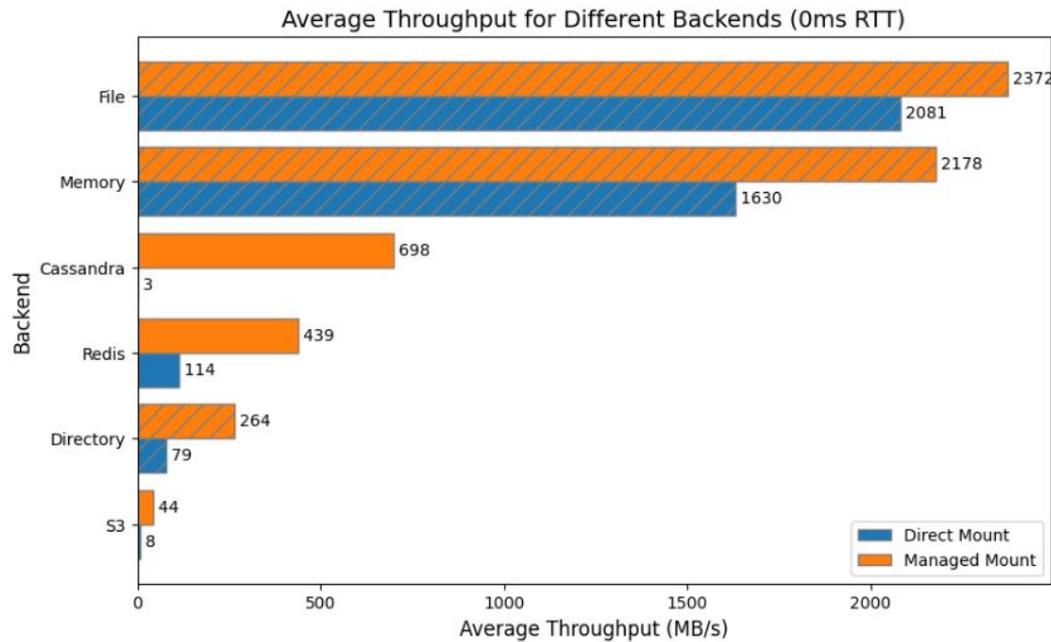


Figure 28: Average throughput for memory, file, directory, Redis, S3 and ScyllaDB backends for direct and managed mounts (0ms RTT)

Throughput

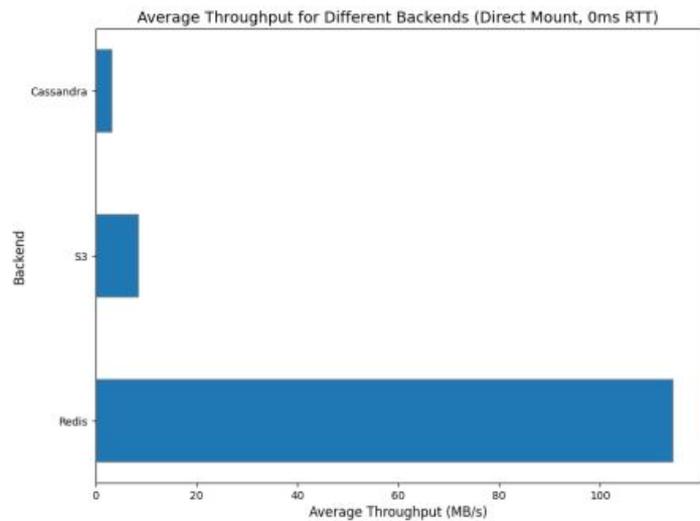


Figure 29: Average throughput for Redis, S3 and ScyllaDB backends for direct mounts (0ms RTT)

Throughput

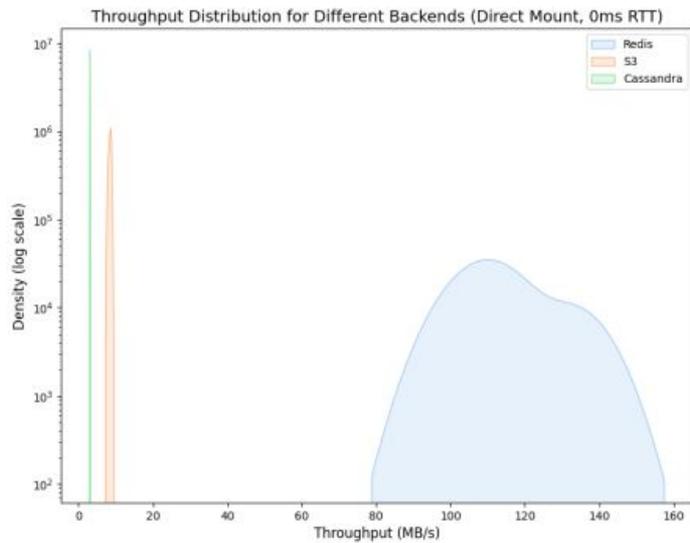


Figure 30: Kernel density estimation (with logarithmic Y axis) for the throughput distribution for Redis, S3 and ScyllaDB for direct mounts (0ms RTT)

Throughput

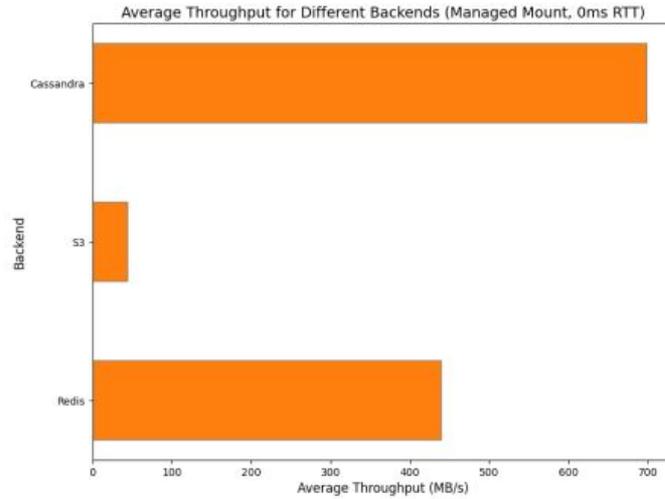


Figure 31: Average throughput for Redis, S3 and ScyllaDB backends for managed mounts (0ms RTT)

Throughput

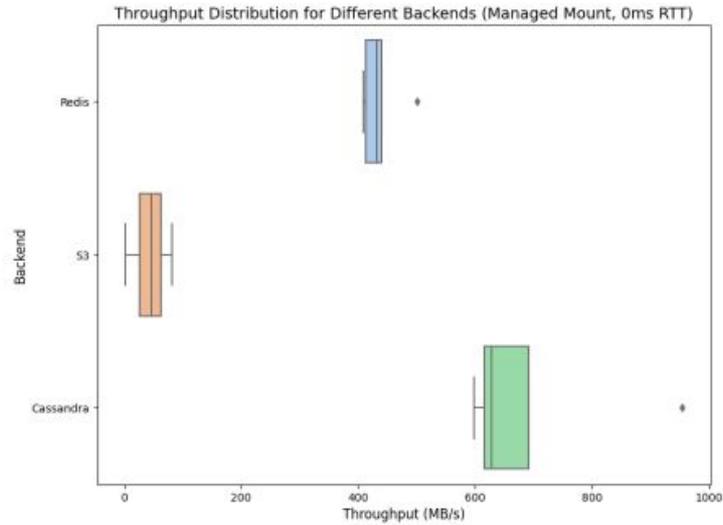


Figure 32: Box plot for the throughput distribution for Redis, S3 and ScyllaDB for managed mounts (0ms RTT)

Throughput

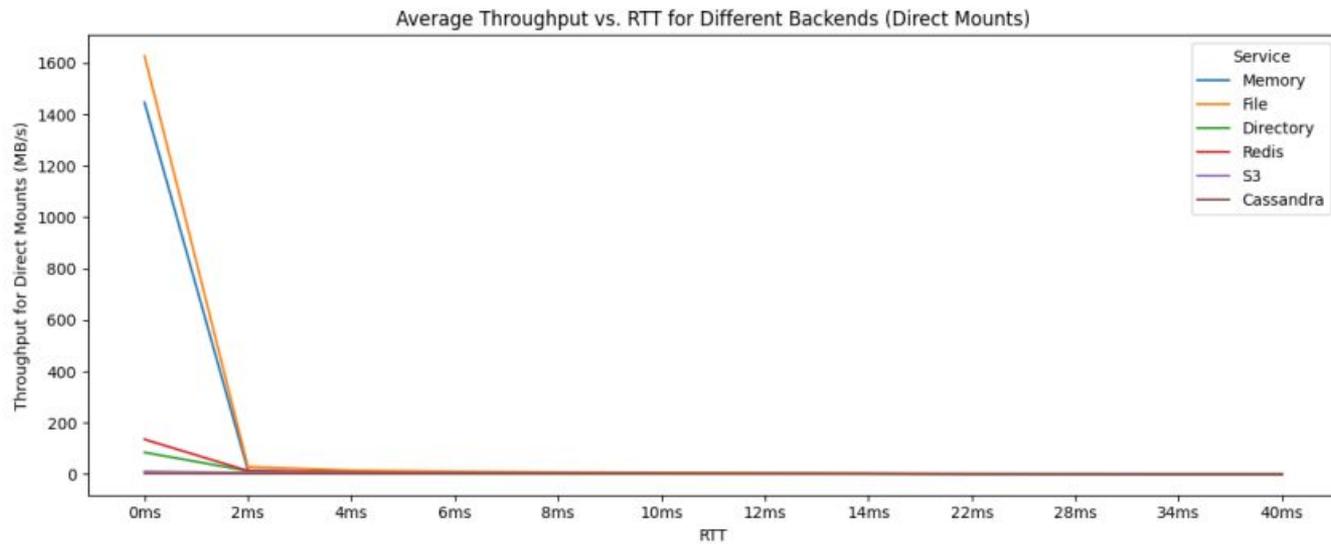


Figure 33: Average throughput for memory, file, directory, Redis, S3 and ScyllaDB backends for direct mounts by RTT

Throughput

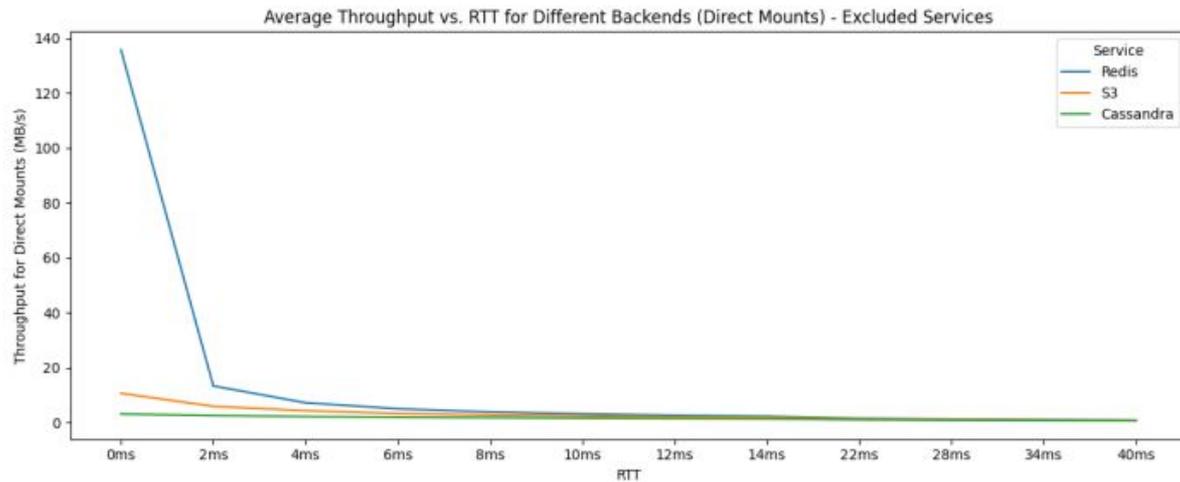


Figure 34: Average throughput for Redis, S3 and ScyllaDB backends for direct mounts by RTT

Throughput

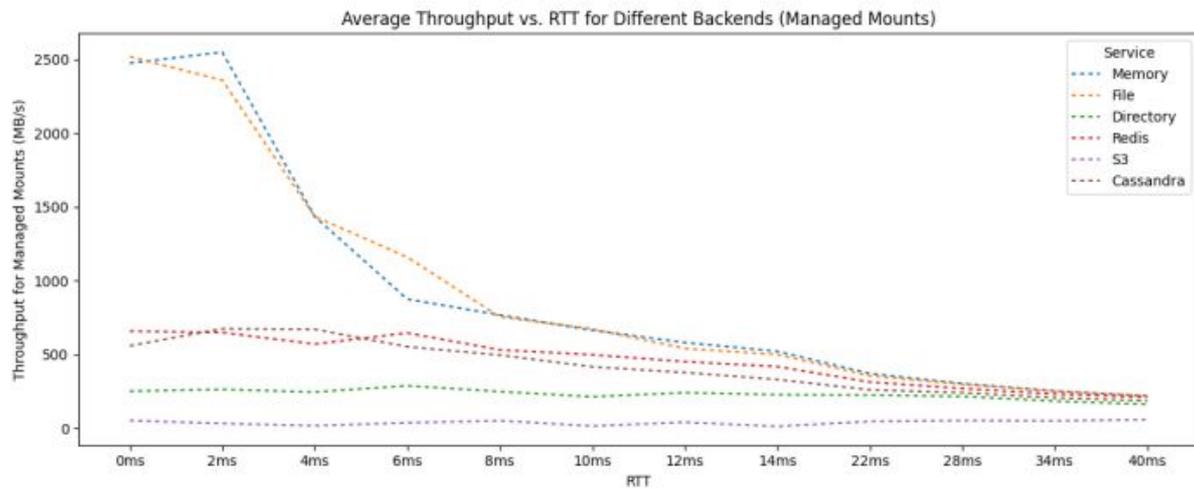


Figure 35: Average throughput for memory, file, directory, Redis, S3 and ScyllaDB backends for managed mounts by RTT

Throughput

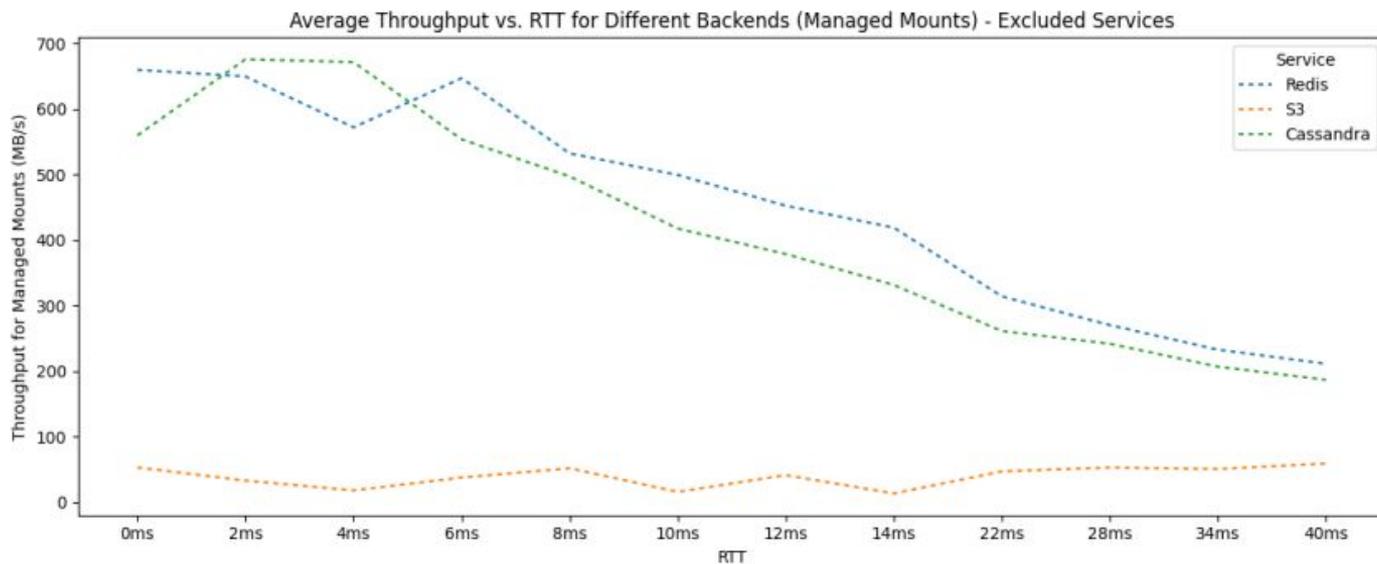


Figure 36: Average throughput for Redis, S3 and ScyllaDB backends for managed mounts by RTT

Throughput

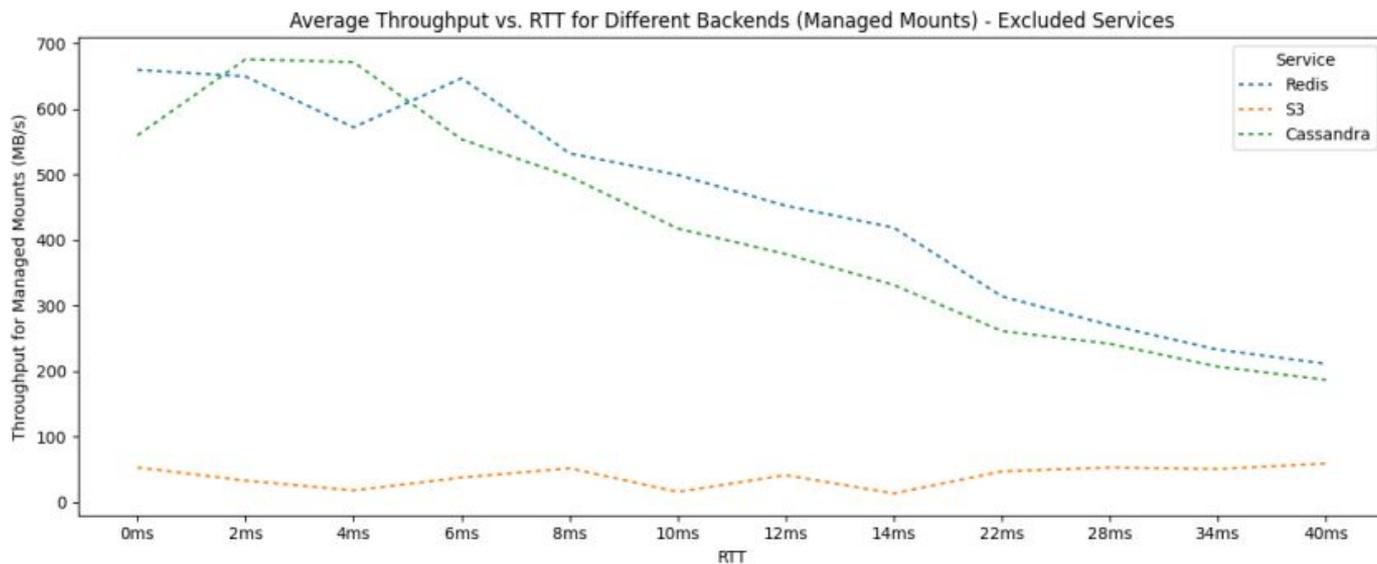


Figure 36: Average throughput for Redis, S3 and ScyllaDB backends for managed mounts by RTT

Throughput

Discussing Backends

Implemented Use Cases



ram-dl

ram-dl

Usage

TL;DR: "Upload" RAM with `ram-ul`, "download" the RAM with `ram-dl`, done!

1. Upload RAM

On a remote (or local) system, first start `ram-ul`. This component exposes a memory region, file or directory as a FRPC server:

```
$ ram-ul --size 4294967296
2023/06/30 14:52:12 Listening on :1337
```

2. Download RAM

On your local system, start `ram-dl`. This will mount the remote system's exposed memory region, file or directory using FRPC and r3map as swap space, and amount it as soon as you interrupt the app:

```
$ sudo modprobe nbd
$ sudo ram-dl --raddr localhost:1337
2023/06/30 14:54:22 Connected to localhost:1337
2023/06/30 14:54:22 Ready on /dev/nbd0
```

This should give you an extra 4GB of local memory/swap space, without using up significant local memory (or disk space):

```
# Before
$ free -h

```

	total	used	free	shared	buff/cache	available
Mem:	30Gi	7.9Gi	6.5Gi	721Mi	16Gi	21Gi
Swap:	8.0Gi	0B	8.0Gi			

```
# After
$ free -h

```

	total	used	free	shared	buff/cache	available
Mem:	30Gi	7.9Gi	6.5Gi	717Mi	16Gi	21Gi
Swap:	11Gi	0B	11Gi			

 **That's it!** We hope you have fun using `ram-dl`, and if you're interested in more like this, be sure to check out [r3map!](#)

ram-dl



tapisk

tapisk

```
1 func (b *TapeBackend) ReadAt(p []byte, off int64) (n int, err error) {
2     // Calculating the block for the offset
3     block := uint64(off) / b.blocksize
4
5     // Getting the physical record on the tape from the index
6     location, err := b.index.GetLocation(block)
7     // ...
8
9     // Creating the seek operation
10    mtop := &ioctl.Mtop{}
11    mtop.SetOp(ioctl.MTSEEK)
12    mtop.SetCount(location)
13
14    // Seeking to the record
15    syscall.Syscall(
16        syscall.SYS_IOCTL,
17        drive.Fd(),
18        ioctl.MTIOCTOP,
19        uintptr(unsafe.Pointer(mtop)),
20    )
21    // ...
22
23    // Reading the chunk from the tape into memory
24    return b.drive.Read(p)
25 }
```

tapisk

Future Use Cases

Improving Cloud Storage Clients

Universal Database, Media and Asset Streaming

Universal App State Mounts and Migrations

Conclusion

Thanks!



r3map

Remote mmap: High-performance remote memory region mounts and migrations in user space.

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Overview

r3map is a library that simplifies working with remote memory regions and migrating them between hosts.

It can ...

- Create a virtual `[]byte` or a virtual file that transparently downloads remote chunks only when they are accessed: By providing multiple frontends (such as a memory region and a file/path) for accessing or migrating a resource, integrating remote memory into existing applications is possible with little to no changes, and fully language-independent.
- mmap any local or remote resource instead of just files: By exposing a simple backend interface and being fully transport-independent, r3map makes it possible to map resources such as a **S3 bucket**, **Cassandra** or **Redis database**, or even a tape drive into a memory region efficiently, as well as migrating it over an RPC framework of your choice, such as gRPC.
- Enable live migration features for any hypervisor or application: r3map implements the APIs which allow for zero-downtime live migration of virtual machines, but makes them generic so that they can be used for any memory region, bringing live migration abilities to almost any hypervisor or application with minimal changes and overhead.
- Overcome the performance issues typically associated with remote memory: Despite being in user space, r3map manages (on a typical desktop system) to achieve very high throughput (up to 3 GB/s) with minimal access latencies (~100µs) and short initialization times (~12ms).
- Adapt to challenging network environments: By implementing various optimizations such as background pull and push, two-phase protocols for migrations and concurrent device initialization, r3map can be deployed not only in low-latency, high-throughput local datacenter networks but also in more constrained networks like the public internet.

[pointfx/r3map](#)

High-performance remote memory region mounts and migrations in user space.

Go ★ 21

Efficient Synchronization of Linux Memory Regions over a Network: A Comparative Study and Implementation

Bachelor's thesis by Felicitas Pöjtinger.

University: Hochschule der Medien Stuttgart

Course of Study: Media Informatics

Date: 2023-08-03

Academic Degree: Bachelor of Science

Primary Supervisor: Prof. Dr. Martin Goik

Secondary Supervisor: M.Sc. Philip Betzler

[Deliverance](#) [Failing](#)

Abstract

Current solutions for access, synchronization and migration of resources over a network are characterized by application-specific protocols and interfaces, which result in fragmentation and barriers to adoption. This thesis aims to address these issues by presenting a universal approach that enables direct operation on a memory region, circumventing the need for custom-built solutions. Various methods to achieve this are evaluated on parameters such as implementation overhead, initialization time, latency and throughput, and an outline of each method's architectural constraints and optimizations is provided. The proposed solution is suitable for both LAN and WAN environments, thanks to a novel approach based on block devices in user space with background push and pull mechanisms. It offers a unified API that enables mounting and migration of nearly any state over a network with minimal changes to existing applications. Illustrations of real-world use cases, configurations and backends are provided, together with a production-ready reference implementation of the full mount and migration APIs via the open-source r3map (remote mmap) library.



Felicitas Pöjtinger, 2023

Efficient Synchronization of Linux Memory Regions over a Network
A Comparative Study and Implementation



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Organizations



Hi there 🙋

I'm Felicitas 🙋, an experienced software engineer with a passion for systems programming in Go, all things DevOps and the modern web 🌐.

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